

**COMMUNICATIONS SYSTEM AND ASSOCIATED METHODS
WITH OUT-OF-BAND CONTROL**

Related Application

5 The present application is based upon
provisional patent application serial no. 60/112,379
filed on December 14, 1998, and the entire disclosure
of which is incorporated herein by reference.

10 **Field of the Invention**

The invention relates to communications
systems and methods, and, more particularly, to digital
communications systems and associated methods over
parallel communications channels.

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Background of the Invention

Digital communications are widely used for
the transmission of voice, data and video information.
Such transmission can extend over large geographical
20 distances, between components within a personal
computer, or only between adjacent circuit portions on
an integrated circuit. Certain such communications
applications benefit from or require the conversion of
serial data into parallel data for simultaneous
25 transmission over parallel communications channels, or
more generically, from M'ary symbols to N'ary symbols.
At the receiving end, the parallel data is desirably
converted back into the serial data, and with the bits
or symbols in the correct order to avoid data errors.

Unfortunately, the demand for greater data transmission volumes and at ever higher speeds, may result in skew at the receiver. In other words, the parallel communications channels may introduce
5 different delays to the parallel symbol strings they carry. Because of skew, the parallel symbol strings at the receiver can then no longer be simply reassembled into the starting data.

The skew problem with parallel communications
10 channels has been addressed in a number of ways. For example, U.S. Patent No. 4,677,618 to Haas et al. recognized the dispersion introduced by wavelength division multiplexed communications channels over optical fiber. This patent discloses determining the
15 relative delays between the channels based upon detecting two bits in a given byte of data. The relative times of arrival of the remaining bits in a byte are predetermined using the relative delay between the two detected bits and the known frequency-related
20 dispersion characteristics of the transmission medium. Certain bits in each received byte may then be delayed using clock delay lines or registers, thereby accounting for skew.

Along similar lines, U.S. Patent No.
25 5,157,530 to Loeb et al. also determines and accounts for skew imparted by dispersion in fiber optic wavelength division multiplexing. Relative delays are used to control adjustable delay devices in each channel.

30 U.S. Patent No. 5,408,473 to Hutchinson et al. is directed to a technique for synchronizing run-length-limited data transmitted over parallel communications channels. Block boundary synchronization is established during connection
35 initialization by using a property of a required HALT code to detect block boundaries received in each channel. Skew compensation is effected by comparing

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the times of detection of the block boundaries in the two channels, and appropriately controlling a variable delay in at least one of the channels. If there is a subsequent loss of synchronization, detected

5 transmission errors will eventually result in connection reinitialization and reestablishment of synchronization. Unfortunately, the transmission of the fixed HALT code to detect boundaries may result in false boundary detection. Moreover, since
10 synchronization is not continuously maintained, the technique may be impractical for higher data rates.

U.S. Patent No. 5,793,770 to St. John et al. is directed to a high-performance parallel interface (HIPPI) to a synchronous optical network (SONET)
15 gateway, and wherein electronic logic circuitry formats data and overhead signals into a data frame for transmission over a fiber optic channel. Stripe skew adjustment is based upon SONET framing, and, as such, the circuitry is relatively complicated, comprising as
20 many as 20,000 logic gates, for example.

The difficulty with skew caused by parallel communications channels is also an important issue to be addressed in communications channels between integrated circuit devices. For example, higher
25 transmission speeds increase the sensitivity to skew, as there is a smaller time window to correctly identify a received bit and have it properly align with bits received on the other parallel communications channels. To provide a higher aggregate transmission rate, the
30 number of parallel communications channels can be increased, without increasing the speed of any given communications channel. However, this may result in significant costs for the additional communications channels. Moreover, for communications between
35 integrated circuits, increasing the number of communications channels increases the number of pins needed for connecting the IC. The number of pins and

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additional packaging complexity may significantly increase the costs of such approaches.

For communications channels between physical layer devices (PLDs) or PHY devices, and logical link devices (LLDs), typical interfaces are asymmetrical and the devices are operated in a push-pull configuration. Because of the asymmetry, relatively expensive memory is required on the PLD since it is polled by the LLD, such as an asynchronous transfer mode (ATM) device. Further developments and improvements in the communications interface between a PLD and LLD are also hampered by the skew difficulty described above as a result of higher bit rates over limited parallel communications channels.

Summary of the Invention

In view of the foregoing background, it is therefore an object of the present invention to provide a communications system and associated methods which overcome the drawback of the prior art, particularly relating to more efficient control of the PLD device by the LLD.

These and other objects, features and advantages in accordance with the present invention are provided by a communications system comprising a PLD having a PLD send interface which, in turn, includes PLD parallel information outputs and at least one PLD control output; and an LLD comprising an LLD receive interface which, in turn, includes LLD parallel information inputs and at least one LLD control input. The communications system also preferably includes first parallel communications channels connecting the PLD information outputs to respective LLD information inputs, and at least one second communications channel connecting the at least one PLD control output to the at least one LLD control input so that control signals are sent from the PLD to the LLD out-of-band from

information signals. Accordingly, control speed is enhanced, and information throughput efficiency is not compromised.

5 The LLD receive interface further preferably includes at least one LLD status output, and the PLD send interface preferably further includes at least one PLD status input. At least one third communications channel thus preferably connects the at least one LLD status output to the at least one PLD status input.

10 The PLD may further include a PLD receive interface including PLD parallel information inputs and at least one PLD control input, and the LLD may further comprises an LLD send interface including LLD parallel information outputs and at least one LLD control
15 output. In this embodiment, fourth parallel communications channels connect the LLD information outputs to respective PLD information inputs. In addition, at least one fifth communications channel preferably connects the at least one LLD control output
20 with the at least one PLD control input. In accordance with this aspect of the invention, the PLD and LLD are operable in a push-push configuration. The push-push configuration overcomes a number of difficulties of conventional push-pull configurations.

25 Another advantageous feature of the invention is that the interfaces may be symmetric. In other words, the PLD send interface and the LLD send interface may be substantially identical, and the PLD receive interface and the LLD receive interface may be
30 substantially identical to thereby define symmetrical interfaces. Symmetrical interfaces may simplify design and manufacturing and offer other advantages including loop-back capabilities and additional device packaging options.

35 The PLD receive interface may further includes at least one PLD status output, the LLD send interface may further include at least one LLD status

input, and the system may also include at least one sixth communications channel connecting the at least one PLD status output to the at least one LLD status input.

5 The LLD may comprise an asynchronous transfer mode (ATM) device, for example. In addition, the PLD may be a synchronous optical network (SONET) device or a synchronous digital hierarchy (SDH) device.

Yet another aspect of the invention is that
10 the pin count of the PLD and LLD may be kept manageable by using higher speed parallel communications channels while accounting for skew. In particular, the PLD send interface may comprise a string-based framing coder for determining and appending a string-based framing code
15 to each information symbol string of information symbol strings to be transmitted in parallel over respective first parallel communications channels, with each string-based framing code being based upon at least some of the information symbols in the respective
20 information symbol string. Also, the LLD receive interface may comprise a deskewer for aligning received parallel information symbol strings based upon the string-based framing codes. The information symbols may be binary bits, and the string-based framing codes
25 may be CRC codes, for example.

The deskewer may comprises a framer for framing information symbol strings based upon the respective string-based framing codes, and an aligner for aligning framed information symbol strings relative
30 to one another and based upon the string-based framing codes. The aligner, in turn, may comprise at least one first-in-first-out (FIFO) device connected to the framer for buffering framed information bit strings. The aligner may also include a FIFO controller for
35 aligning framed information symbol strings during at least one of a writing and a reading phase of the at

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The coding and deskewing may be provided in the reverse direction on information and control

A method aspect of the invention is for communicating including the steps of: sending information signals over first parallel communications channels from the PLD to the LLD, and while sending control signals over at least one second communications channel from the PLD to the LLD so that control signals are sent from the PLD to the LLD out-of-band from information signals. This arrangement permits the PLD to operate without unnecessary throttling from the LLD. The arrangement also overcomes difficulties which would otherwise occur if control signals were included in with the information signals.

The method may also include sending status signals over at least one third communications channel from the LLD to the PLD. The step of sending status signals over at least one third communications channel may comprise the steps of operating a PLD send interface including at least one PLD status input, and

operating an LLD receive interface including at least one LLD status output.

The method may also include the steps of sending information signals over fourth parallel communications channels from the LLD to the PLD, and while sending control signals over at least one fifth communications channel from the PLD to the LLD so that control signals are sent from the PLD to the LLD out-of-band from information signals. The step of sending information signals over fourth parallel communications channels may include the steps of operating an LLD send interface including LLD parallel information outputs; and operating a PLD receive interface including PLD parallel information inputs. The step of sending control signals over at least one fifth communications channel preferably comprises the steps of operating an LLD send interface including at least one LLD control output, and operating a PLD receive interface including at least one PLD control input. In addition, the method may also include the step of sending status signals over at least one sixth communications channel from the PLD to the LLD.

The method in accordance with the present invention permits the PLD and LLD to be operated in a push-push configuration. This overcomes the drawbacks of prior approaches using push-pull configurations. Considered in somewhat simpler terms, the PLD may comprise a PLD send interface and the LLD may comprise an LLD send interface substantially identical to the PLD send interface. Additionally, the PLD may comprise a PLD receive interface and the LLD may comprise an LLD receive interface substantially identical to the PLD receive interface thereby define symmetrical interfaces.

Brief Description of the Drawings

FIG. 1 is a schematic block diagram of a first embodiment of a communications system in accordance with the present invention.

5 FIG. 2 is a schematic bit position diagram from the output of the first device as shown in FIG. 1.

FIG. 3 is a schematic bit position diagram from the input of the second device as shown in FIG. 1 and illustrating skew.

FIG. 4 is a schematic bit position diagram from the FIFO device of the second device as shown in FIG. 1 and illustrating deskewing.

FIG. 5 is a schematic diagram of a framing state machine as may be used in the second device as shown FIG. 1.

FIG. 6 is a schematic block diagram of an optical fiber embodiment of a communications system in accordance with the present invention.

20 FIG. 7 is a schematic block diagram of a radio embodiment of a communications system in accordance with the present invention.

FIG. 8 is a schematic block diagram of an infrared free space embodiment of a communications system in accordance with the present invention.

25 FIG. 9 is a schematic block diagram of another embodiment of a communications system illustrating a bank of lower rate converter electronics and incorporating the deskewing features in accordance with the present invention.

FIG. 10 is a schematic block diagram of an optical fiber embodiment of a communications system as shown in FIG. 9.

35 FIG. 11 is a schematic block diagram of a communications system including a PLD and an LLD connected by parallel communications channels in accordance with the present invention.

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FIG. 12 is a more detailed schematic block diagram of the PLD send interface and LLD receive interface as shown in FIG. 11.

FIG. 13 is a schematic diagram of binning of
5 the data and control bits for the PNG interface example in accordance with the present invention.

FIG. 14 is a schematic block diagram of an example of a deskewing algorithm in accordance with the present invention.

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Detailed Description of the Preferred Embodiments

The present invention will now be described more fully hereinafter with reference to the accompanying drawings, in which preferred embodiments
15 of the invention are shown. This invention may, however, be embodied in many different forms and should not be construed as limited to the embodiments set forth herein. Rather, these embodiments are provided so that this disclosure will be thorough and complete,
20 and will fully convey the scope of the invention to those skilled in the art. Like numbers refer to like elements throughout, and prime and multiple prime notation is used in alternate embodiments to refer to similar elements.

25 Referring initially to FIGS. 1-5, an embodiment of a communications system 20 in accordance with the present invention is first described. The communications system 20 illustratively includes a first device 22 and a second device 24 connected by
30 parallel communications channels. In the illustrated embodiment, the communications channels are provided by electrical conductors or wires 25-29, although other transmission media may be used to establish or define the parallel communications channels as will be readily
35 appreciated by those skilled in the art. Also in the illustrated embodiment four wires 25-28 are connected

to carry information bits, while the fifth wire 29 carries a clock signal from the clock 42. The communications channel for the clock signal is not needed in all embodiments, as the clock signal can typically be recovered if the received information bits have a sufficient number of transitions as will also be appreciated by those skilled in the art.

As explained above in the Background of the Invention, parallel communications channels may present a skew problem especially where the bit rate is relatively high or the distance is relatively long. For example, for an 800 Mbs rate over electrical parallel conductors, skew may limit separation distances to two inches or less.

For clarity of explanation, the following description will be with reference to transmitting binary information elements or information bit strings. In other words, the term "information bit string" will be used, although those of skill in the art will understand that symbols other than binary one's and zero's can also be used in accordance with the present invention. For example, a three-level information symbol may also be used and benefit from the deskewing concepts described herein.

The first device 22 illustratively includes a string-based framing coder 32 for determining and appending a string-based framing code to each information bit string of information bit strings to be transmitted in parallel over respective parallel communications channels. "Appending" is meant to cover both prepending and postpending, although those skilled in the art will recognize that postpending may be preferred, since prepending may require more buffer memory.

A scrambler 34 is connected upstream from the string-based framing coder 32. The scrambler 34 may be

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desirable to avoid long strings of null values which could hinder clock recovery as will be appreciated by those skilled in the art. In other embodiments, the scrambler 34 may be connected downstream from the

5 string-based framing coder 32. Such an optional scrambler may be a self-synchronizing scrambler, such as an X^{43} scrambler as will be appreciated by those skilled in the art.

An M'ary-to-N'ary mapper or converter 36 is
10 illustratively connected upstream of the scrambler 34. The M'ary-to-N'ary converter 36 converts the incoming Mbits to N parallel information bit strings for subsequent transmission over the parallel communications channels. The M'ary-to-N'ary converter
15 36 is conventional and requires no further discussion herein.

Returning again to the string-based framing coder 32, this illustratively includes a string-based code generator 37 for generating each string-based
20 framing code based upon at least some of the information bits in the respective information bit string. A multiplexer 38 appends the string-based codes to the respective information bit strings as will be appreciated by those skilled in the art. An
25 electrical-to-medium converter 41 is connected between the output of the multiplexer 38 and the communications channels provided by the wires 25-29. In this illustrated embodiment, the electrical-to-medium converter 41 may be provided by suitable electrical
30 driver circuitry as will be appreciated by those skilled in the art. In other embodiments, the electrical-to-medium converter 41 may connect to other transmission media.

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The second device 24 preferably includes a deskewer 45 for aligning received parallel information bit strings based upon the string-based framing codes. The string-based framing codes and their use to deskew received information bit strings permit the information bits to be transmitted at high rates and/or over relatively long distances.

In one preferred embodiment, the string-based coder 32 comprises a cyclic redundancy checking (CRC) coder for determining and appending CRC codes to respective information bit strings. Thus, the deskewer may comprise a CRC framer for framing the information bit strings based upon the CRC codes. Of course, the second device 24 may also include an error detection and correction circuit 47 using the CRC codes. Each CRC code may be one of a CRC-4 to CRC-32 code, for example. For an information bit string of 1024 bits, for example, a CRC-8 code may be sufficient to ensure quick and accurate framing.

The string-based code may also include other bits in addition to those specifically based on the information bit string, such as the CRC code bits, for example. Some bits may be assigned as counting or identifying bits to be used when the expected delay or skew was greater than a single frame as will be appreciated by those skilled in the art. Of course, other bits could be assigned for other purposes as well.

A particular advantage of the CRC coding is that a straightforward hardware implementation can be achieved with a relatively small number of logic gates as will be appreciated by those skilled in the art. CRC codes are also resistant to false framing while adding relatively little overhead to the information bit strings. Fixed framing in contrast, would likely experience considerable false framing for a similar

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number of code bits. If the number of fixed framing bits were increased to reduce false framing, the overhead may be considerable. Yet another advantage of CRC codes is that they may also be used for error
5 detection and correction as they are conventionally used. Accordingly, the second device 24 may optionally include the illustrated error detect and correction circuit 47, which needs no further discussion herein.

The deskewer 45 may comprise a framer 50 for framing information bit strings based upon the respective string-based framing codes. The illustrated deskewer 45 of the second device 24 also includes an aligner 52 for aligning framed information bit strings relative to one another and based upon the string-based framing codes. The aligner 52 may, in turn, advantageously comprise at least one first-in-first-out (FIFO) device 53 connected to the framer 50 for buffering framed information bit strings as shown in the illustrated embodiment. The aligner 52 also illustratively includes a FIFO controller 55 for aligning framed information bit strings during at least one of a writing and a reading phase of the at least one FIFO device and based upon the string-based framing codes. The term "FIFO device" is used herein to include a FIFO, a shift register, and any other type of ordered storage element as will be appreciated by those skilled in the art.

30 All of the information bit strings may have a same number of bits in some embodiments to simplify the system implementation. In other embodiments, the bit strings could have different lengths as would be appreciated by those skilled in the art.

Turning now to the front end of the second device 24, a sampler 56 is connected upstream from the deskewer 45. The sampler 56 samples the received bit

string based upon the clocking pulses as will be appreciated by those skilled in the art. Ideally the sampler 56 samples the bit string at a bit midpoint. The clock signal for the sampler 56 may come from the
5 clock receiver 57 or from the recovered clock 58, the operation of both of which will be appreciated by those skilled in the art.

A medium-to-electrical converter 61 is connected between the sampler 56 and the communications
10 channels as provided by the wires 25-29. Of course, other types of converters can be used for different transmission media.

The deskewer 45 also illustratively includes a descrambler 46 for descrambling the information bit
15 strings, such as to facilitate clock recovery at the second device 24. The descrambler 46 is illustratively connected between the framer 50 and the FIFO device 53. In other embodiments, the descrambler 46 may be connected downstream from the FIFO device 53 as will be
20 appreciated by those skilled in the art. Of course, in other embodiments, the descrambler 46 and the scrambler 34 may not be needed at all.

Referring now more particularly to FIGS. 2-4, the deskewing in accordance with the invention is
25 described with reference to a simplified example. The table 62 in FIG. 2 illustrates the alignment of some of the information bits A-P and some of the CRC bits C11-C42. This is the proper alignment that would typically be produced at the output of the first device 22 or at
30 a relatively short distance therefrom as will be appreciated by those skilled in the art.

As shown in the table 63 of FIG. 3, the second information bit string from the top is out of alignment with the other information bit strings.

Accordingly, the information bit string that would be recovered without deskewing would be A, Z, C, . . . P. In other words, the information bit string would be incorrect.

Now, as shown in the table 64 of FIG. 4, the deskewing of the present invention re-aligns the frames that may have been misaligned due to skew. Accordingly, the correct information bit string, A, B, . . . P, is produced at the output. The communications system 20 and associated deskewing method using the string-based framing codes advantageously and efficiently removes or accounts for the skew. This permits higher bit rates and/or longer transmission distances. The higher bit rates may permit a reduction of pin count for communication between integrated circuit chips. As the cost for additional pins and packaging complexity may be relatively high, the present invention also permits lower cost communications ICs having an aggregate communication rate that is still relatively high as will be appreciated by those skilled in the art. As will also be appreciated by those skilled in the art, although only one direction of communication has been illustrated, other embodiments of the communications system 10 may include circuitry to implement a reverse direction of communications. In other words, the present invention is also readily applicable to full-duplex communications systems as well. In addition, multiple receiving devices may be connected to one or more transmitting devices as will also be appreciated by those skilled in the art.

Referring now more particularly to FIG. 5, some additional framing or synchronization concepts are now explained. In particular, the illustrated state framing machine 70 has three states: a pre-sync state 35 71, a hunt state 73 and a sync state 72. Transition

between the hunt and pre-sync states is determined based upon a correct or incorrect string-based framing code. The machine 70 changes from the pre-sync state 71 to the sync state 72 if X consecutive correct codes are determined. The machine 70 transitions from the sync state 72 to the hunt state 73 if there are Y consecutive incorrect codes determined. The state framing machine 70 is very similar to state framing machines used in other known data synchronizing applications as will be appreciated by those skilled in the art.

The fundamental string-based framing coding and associated deskewing concepts of the present invention have initially been explained with reference to parallel communications channels provided over wires 25-59. In other words, the communication system 20 operates over a wireline transmission medium. Another wireline application would include operation over a data bus, such as a PCI bus, for example. A typical PCI bus is limited to a relatively low frequency of about 60-70 MHZ. Accordingly, for greater information throughput wider buses are needed. The present invention overcomes this difficulty and can allow a PCI bus to operate at a faster clock speed without additional bus width. Other wireline transmission media include twisted copper pairs, and coaxial cables, for example, as will be appreciated by those skilled in the art.

Turning now additionally to FIGS. 6-8, several alternate embodiments with respect to the transmission media are now described. For the communications system 20' shown in FIG. 6, the communications medium is illustratively provided by two optical fibers 75. The first device 22' includes the string-based framing coder 32', and the second device

24' includes the deskewer 45'. The communications system 20' may also include other components as shown in FIG. 1 and described above.

A first wireless communications system embodiment 20" is explained with particular reference to FIG. 7. In this embodiment, radio transmitters and receivers 76, 77, respectively, and free space provide the communications channels over a wireless medium. The other components are indicated with double prime notation and are similar to those described above.

FIG. 8 illustrates a second wireless communications system 20"' wherein infrared transmitters and infrared detectors 81, 82, respectively, provide the parallel communications channels over free space as will be appreciated by those skilled in the art. Those of skill in the art will also appreciate other more specific implementations and applications of the general communications systems 20, 20', 20" and 20"' described herein.

Other advantageous features and implementations of the present invention are now described with additional reference to FIGS. 9 and 10. More particularly, as shown in FIG. 9, the invention is also directed to a communications system 120 including a first device 122 comprising a plurality of electrical-to-transmission medium converters 141, and a second device 124 comprising a plurality of transmission medium-to-electrical converters 161. The electrical-to-transmission medium converters 141 are connected to respective ones of the transmission medium-to-electrical converters 161 via at least one transmission medium and defining parallel communications channels between the first and second

devices. In the illustrated embodiment, the transmission medium is provided by the wireline 125.

In other embodiments, wireless and optical transmission media may be used. A radio wireless
5 medium is schematically indicated by antennas 113, 114, and an infrared or free space optical medium is indicated by source 115 and detector 116.

The first device 122 comprises a string-based framing coder 132 for determining and appending a
10 string-based framing code to each information bit string of information bit strings to be transmitted in parallel over respective parallel communications channels, each string-based framing code being based upon at least some of the information bits in the
15 respective information bit string as described above with respect to the string-based framing coder 32 as described above with reference to FIG. 1. The string-based coder 132 may include a string-based code generator and multiplexer as also described above with
20 respect to FIG. 1.

The second device 124 comprises a deskewer 145 for aligning received information bit strings based upon the string-based framing codes. The deskewer 145 may include the components and/or equivalents as
25 described above for the deskewer 45 shown in the communications system 20 of FIG. 1. Accordingly, lower rate converters 141, 161 can be used in the communication system 120 at significant costs savings and while providing a desired relatively high overall
30 information throughput rate. The deskewing features account for any skewing that may occur through the parallel communications channels.

For longer distances, the first device 122 may further comprise a multiplexer 110 for multiplexing

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signals from the plurality of electrical-to-transmission medium converters 141 along a common transmission medium, such as the wireline medium 125. In these embodiments, the second device 124 also
5 includes a demultiplexer 112 connected to its transmission medium-to-electrical converters 161 as will be appreciated by those skilled in the art. The medium multiplexing and demultiplexing may reduce the cost for the overall communications system 120 for
10 relatively large distances between the first and second devices 122, 124 as will also be appreciated by those skilled in the art.

For other embodiments, the multiplexer 110 and demultiplexer 112 may not be needed. Those of
15 skill in the art will readily be able to determine the cost tradeoffs to implement the communications system 120 either with or without the mux/demux.

One particularly advantageous use of the tradeoff in converter speed versus number of parallel
20 communication channels is for optical implementations, such as for those in accordance with the synchronous optical network (SONET) and/or synchronous digital hierarchy (SDH) standards. In particular, wavelength division multiplexing (WDM) and dense wavelength
25 division multiplexing (DWDM) are highly developed and permit a relatively large number of communications channels to be established at different optical wavelengths as will be appreciated by those skilled in the art. In the illustrated communications system
30 120', a larger number of lower rate OC-X converters 141', 161' may be used in place of a lesser number of higher rate converters. The aggregate rate across the fiber 175 can be relatively high, that is, N times the OC-X rate. For example, an OC-192 converter may be 100
35 times the cost of a OC-48 converter. Accordingly, the

communications system 120' may be less expensive than comparable aggregate transmission rate systems using higher speed optical converters.

5 The different wavelengths multiplexed onto the fiber 175 travel at different speeds through the fiber. In the past, these different speeds have been measured and a fixed offsets added to account for skew. Unfortunately, such approaches may not account of changes in skew as may be caused by fiber aging and/or
10 temperature cycling, for example. Those of skill in the art will appreciate the advantages in terms of efficiency and simplicity provided by the string-based framing coder 132' and deskewer 145' and associated methods in accordance with the invention for the
15 optical communications system 120'.

Another communications system 200 in accordance with the invention is now described with particular reference to FIGS. 11 and 12. The illustrated communications system 200 includes a
20 physical layer device (PLD) 201 and a logical link device (LLD) 202 connected thereto. The PLD 201 includes a PLD send interface 203 which, in turn, includes PLD parallel information outputs 205-208 and at least one PLD control output 211.

25 The LLD 202 comprises an LLD receive interface 204 which, in turn, includes LLD parallel information inputs 215-218 and at least one LLD control input 221. The communications system 200 also includes first parallel communications channels 225 connecting
30 the PLD information outputs 205-208 to respective LLD information inputs 215-218. A second communications channel 226 connects the illustrated single PLD control output 211 to the LLD control input 221 so that control

signals are sent from the PLD to the LLD out-of-band from information signals. Accordingly, control speed is enhanced, and information throughput efficiency is not compromised. The number of first and second
5 communications channels 225 and 226 can be different in different embodiments, as will be appreciated by those skilled in the art.

The LLD receive interface 204 further illustratively includes an LLD status output 222, and
10 the PLD send interface 203 includes a PLD status input 212. A third communications channel 227 thus connects the LLD status output 222 to the PLD status input 212.

The PLD 201 also includes a PLD receive interface 230 including PLD parallel information inputs
15 231-234 and a PLD control input 235. The LLD 202 further comprises an LLD send interface 240 including LLD parallel information outputs 241-244 and an LLD control output 245. In the illustrated embodiment, fourth parallel communications channels 250 connect the
20 LLD information outputs 241-244 to respective PLD information inputs 231-234. In addition, a fifth communications channel 251 connects the LLD control output 245 with the PLD control input 235.

The PLD receive interface 230 may further
25 include a PLD status output 236, the LLD send interface 240 may further include an LLD status input 246, and the communications system 200 may also include a sixth communications channel 252 connecting the PLD status output 236 to the LLD status input 246. In accordance
30 with this aspect of the invention, the PLD 201 and LLD 202 are operable in a push-push configuration.

The push-push configuration overcomes a number of difficulties of conventional push-pull configurations. The PLD 201 typically includes on-chip memory 260 which is relatively expensive. Although the
5 LLD 202 typically requires more memory, the memory 261 is less expensive since it is typically provided off-chip as will be appreciated by those skilled in the art. The push-push configuration relaxes the need for relatively expensive on-chip memory 260 for the PLD
10 201, since the LLD 202 does not throttle the PLD 201.

Another advantageous feature of the invention is that the interfaces 203, 204, 230 and 240 may be made symmetric. In other words, the PLD send interface 203 and the LLD send interface 240 may be substantially
15 identical, and the PLD receive interface 230 and the LLD receive interface 204 may be substantially identical to thereby define symmetrical interfaces. Symmetrical interfaces may simplify design and manufacturing and offer other advantages including
20 loop-back capabilities as indicated by the dotted loopback path 263 as illustrated between the PLD send interface 203 and the PLD receiver interface 230 as will be appreciated by those skilled in the art.

The symmetry also permits the LLD 202 to be
25 provided in two separate integrated circuit packages functionally divided as indicated by the dashed horizontal line 264. Since the number of pins can be relatively large and account for a significant portion of the cost, the symmetry and ability to provide two
30 separate ICs is a significant advantage of the present invention.

The LLD 202 may comprise an asynchronous transfer mode (ATM) device, for example, as will be appreciated by those skilled in the art. In addition,

the PLD may be a synchronous optical network (SONET) device or a synchronous digital hierarchy (SDH) device as will also be appreciated by those skilled in the art. The LLD 202 may also be an SDL or HDLC device as
5 will also be appreciated by those skilled in the art.

Yet another aspect of the communications system 200 is that the pin count of the PLD 201 and LLD 202 may be kept manageable by using higher speed parallel communications channels while accounting for
10 skew. In particular, as shown with particular reference to FIG. 11, the PLD send interface 203 may comprise a string-based framing coder 32" for determining and appending a string-based framing code to each information bit string of information bit
15 strings to be transmitted in parallel over respective first parallel communications channels, with each string-based framing code being based upon at least some of the information bits in the respective information bit string. Also, the LLD receive
20 interface 204 may also comprise a deskewer 45" for aligning received parallel information bit strings based upon the string-based framing codes.

The string-based framing codes may be CRC codes, for example. The illustrated string-based coder
25 32" includes a string-based code generator 37" and multiplexer 38" as described above. The scrambler may be added, but is not shown in the illustrated embodiment. An M'ary-to-N'ary converter 36" is upstream from the string-based framing coder 32".

30 The deskewer 45" illustratively includes a framer 50" for framing information bit strings based upon the respective string-based framing codes, and an aligner 52" for aligning framed information bit strings relative to one another and based upon the string-based
35 framing codes. The aligner 52" includes least one

first-in-first-out (FIFO) device 53" connected to the
framer for buffering framed information bit strings;
and a FIFO controller 55" for aligning framed
information bit strings during at least one of a
5 writing and a reading phase of the at least one FIFO
device and based upon the string-based framing codes. A
particular example of an efficient framing algorithm is
provided in the next portion of this description.

The coding and deskewing may be provided in
10 the reverse direction on information and control
signals from the PLD to the LLD. In some embodiments,
the parallel communications channels may be provided
over electrical conductors as shown in the illustrated
embodiment. Of course, other transmission media may
15 also be used.

I. Example Interface

Having already described the general
components, concepts, features and advantages of the
20 present invention, this description now includes a
specific example of an interface entitled the PNG
interface. Of course, this example is meant to further
describe the invention by way of example and should not
be construed to be limiting of the invention.

25 PNG is an interface that allows the transport
of data from one device to another, at very high data
rates, with a minimal number of pins. PNG is
symmetrical, and push/push, and can be used at OC-48
rates extensible to OC-192 and beyond. In addition,
30 PNG is designed to transport both ATM cells and packets
concurrently.

As shown in FIG. 11 it can be see that PNG
includes three kinds of signals: data, control, and
status. Control is sent in the same direction as data,
35 and status is sent in the opposite direction. There
are thus d+c+s bits that comprise an interface per
direction. (Note that unlike Utopia or Utopia like

interfaces, PNG is symmetric, and is thus push/push instead of push/pull.) Table 1 indicates representative or suggested values for d, c and s, for various bit rates.

5

Table 1: Proposed Bit Widths and Rates

	OC-12	OC-48	OC-192
d	4 @ 155 Mb/s	4 @ 622 MB/s	8 @ 1244 MB/s
10 c	1 @ 155 MB/s	1 @ 622 MB/s	2 @ 1244 MB/s
s	1 @ 155 MB/s	1 @ 622 MB/s	2 @ 1244 MB/s

At 622 MB/s and above, controlling skew can be difficult. To ensure alignment between bits, and
15 accomplish framing, each bit has its own framing structure: for every n information bits transmitted, there is a CRC appended. The n information bits, along with the CRC may be referred to as an information frame. The suggested size of n is 1024 bits, and the
20 suggested CRC is a CRC-8 of polynomial x^8+x^2+x+1 . However, this size and CRC can be varied dependent on an analysis to determine maximal probabilistic acquisition times. Other suggested sizes are 512 bits with a CRC-4 or 2048 bits with a CRC-32. The CRC-4
25 requires the least hardware, but may require considerable acquisition time because of high false framing probabilities. The CRC-32 requires more framing circuitry.

In addition to the appended CRC, the data
30 will also be scrambled using a frame synchronous scrambler of sequence length 127, of polynomial x^7+x^6+1 . The scrambler shall be reset to 7'b111_1111 on the most significant bit of the byte following the CRC insertion. Note that the scrambling of information
35 is performed to ensure a uniform distribution of 0's and 1's, to allow for CDR.

It is proposed that a framing state machine similar to HEC delineation be adopted as will be appreciated by those skilled in the art. The framing state machine 70 is shown in FIG. 5, and has already been described.

After frame is found, it is the receiver's responsibility to align the information across multiple bits. In the case of OC-48, the receiver would frame on the 4 d bits and the 1 c bit, and after frame is acquired by all five bits, data and control can be correctly extracted. To accelerate acquisition time, false-positive matches on the CRC can be ignored if frame is not found on all five bits with a limited amount of time (plus or minus a few bits, depending on the allowed skew between bits).

In the case of OC-48 with 4 data bits and 1 control bit operation at 622 MB/s, the data and control would be binned as shown in FIG. 13, forming a relationship of 32 bits for every 8 control bits. This binning of course would be performed after information framing has occurred.

After binning of the control and data has occurred, the control information is used to determine status of the associated data. This information includes Port ID (PID), Group ID (GID), Start of Packet (or cell), and End of Packet (or cell), which byte is the end of the packet. The information is given below in Table 2.

Table 2: Control and Data Fields

8		32	
CONT	PID		DATA
CONT	GID	STATUS	DATA

- The first bit of the control field indicates if the control octet (and its associated data field) is a continuation (i.e. part of) the previous control octet. If CONT is 0, then this is the first octet of a new transmission. Port ID (PID) is 7 bits wide, and is used to identify the port that the data is associated with. This results in up to 128 ports. If more ports are required, the Group ID (GID) field which is 3 bits, is used to extend the range of ports to 1024 (2^{10}).
- 10 The STATUS field is used to convey the rest of the control information. The encoding of the status field is shown in Table 3.

Table 3: Encodings of STATUS

15	0000	NULL	0100	reserved	1000	EOP on previous A	1100	EOP on A
	0001	SOP	0101	reserved	1001	EPO on previous B	1101	EOP on B
	0010	CONT	0110	reserved	1010	EPO on previous C	1110	EOP on C
	0011	ABORT	0111	reserved	1011	EPO on previous D	1111	EOP on D

20

- Because the amount of control information that must be sent for each channel cannot be contained in one 8 bit octet, at least two 8 bit octets must be sent before a port context switch can be performed.
- 25 However, after this, a context switch can be performed on the next 4 byte boundary (implying that the interface is 100% efficient for multiples of 4 byte

packets, as long as the packet is 8 bytes or more).
Note that any size packet, from one byte upwards, can
be sent -- less than 8 bytes reduces the efficiency of
the link.

5 Shown in Table 4 below is an example of null
data being sent, followed by a 32 byte packet.

Table 4: Example of Null Data
Followed by a 32 byte packet

10	0	PID		x
	1	GID	0000	x
15				
	1	GID	0000	x
	0	PID		DATA
	1	GID	0001	DATA
	1	x	0010	DATA
20	1	x	0010	DATA
	1	x	0010	DATA
	1	x	0010	DATA
	1	x	0010	DATA
	1	x	1111	DATA

25 Note that with PID and GID, the number of
channels that can be supported in PNG is 1024. This
could be extended to 2048 if necessary by reducing the
status size to 3 bits and extending the GID to 4 bits.
However, then the data field would have to be utilized
30 to convey information on packets of sizes 1-4 bytes.

 The status field is sent in the opposite
direction of the data and control, and is used to
provide flow control on a per channel basis. The
status field is scrambled, and postpended with a CRC in
35 the same fashion as the control and data. Once the
status information has been extracted, it is

interpreted according to the format shown below in Table 5.

Table 5: Status Information Format

5	7	3	2	4
	PID	GID	R	FSTATUS

As in the Control field, PID represents the Port ID, and GID represents the Group ID, and these combined bits can be used to provide status on up to 1024 ports. R is reserved, and FSTATUS is used to convey the status of the FIFO associated with the port. The R field could be used to grow either the GID or FSTATUS as necessary.

The encoding of the FSTATUS is shown in Table 6 below.

Table 6: Encodings of FSTATUS

	0000	Fifo Empty	0100	1000	1100
20	0001	.	0101	1001	1101
	0010	.	0110	1010	1110
	0011	...	0111	1011	1111	Fifo Full

If only one bit status is required, the MSB bit of the FSTATUS can be employed. The receiving side (defined as the side receiving data and control but sending status) is responsible for determining the status of its FIFOs, and conveying the information to the sending side (defined as the side sending data and control but receiving status). A one bit fifo status can be thought of as similar to TxClav in Utopia interfaces.

The sending side is required to maintain the status of every FIFO on the receiving side. This status is continuously updated whenever a FIFO status changes at the receiver, and in the absence of any such
5 information to convey, the receiver round robins through all active FIFOs to ensure that the sending side is continuously synchronized. In the case of a traditional PHY device as the sender, it does not require status. As such, status information can be
10 ignored by the PHY as a sender.

A logical block diagram of the PNG interface is shown in FIG. 1 and needs no further description.

II. Deskew Algorithm Example

15 Having described the general components, concepts, features and advantages of the present invention, and an interface example this description now includes a specific example of a deskew algorithm. This deskew algorithm is meant to further describe the
20 invention by way of example and should not be construed to be limiting of the invention.

As more bandwidth is needed in a network, customers desire the ability to add bandwidth between two points in an on-demand basis. In a SONET/SDH
25 network this bandwidth is normally provided by sending the next larger concatenation level to the point needing the bandwidth. These networks have major limitations on the timeslot the new concatenated signal can occupy. In accordance with the invention,
30 bandwidth can be allocated/de-allocated at the STS-1/VC-4 level for SONET/SDH signals respectively with no restriction on the STS-1/VC-4 signal used. At the receiving end, a mechanism must exist to align these disjoint streams into one stream without external
35 intervention (software) and minimal hardware (gate count). This example addresses the software

intervention and gate count problem by implementing a completely hardware based deskewing algorithm.

The SONET/SDH SPE deskewing algorithm allows the dynamic alignment of multiple STS-1/VC-4 signals within a user defined group. Accordingly, signals can be added or dropped dynamically from the user defined group without introducing errors in the output signal.

As understood with reference to FIG. 14, the deskewing algorithm 279 takes into account features of a SONET/SDH frame that allow a simple algorithm to align SPEs of different STS-1/VC-4 signals. The algorithm 279 includes four basic components: an input control logic block 280, a write control logic block 281, data storage element block 285, and read control logic block 286.

The input control logic block 280 contains the provisioned information needed to determine which STS-1/VC-4 signals are grouped together to form a Virtual group. This information is also used by other control blocks for deskewing purposes.

The write control logic block 282 contains the write address generation logic and the minimum write address and maximum write address calculation logic for each virtual group. A virtual group includes from two (2) to forty-eight (48) STS-1 signals or two (2) to sixteen (16) VC-4 signals. The write address is synchronized to the J1I marker. This ensures the J1 POH byte is always written into the same location in the data storage memory. The write address always writes 783 (STS-1) or 783*3 (VC-4) bytes between J1 markers. This is independent of increment or decrement operations that have occurred to the incoming pointer values as the signal passed through pointer processor elements in the network. Simultaneously, during each row the maximum and minimum write addresses are calculated for all virtual groups. This information is

